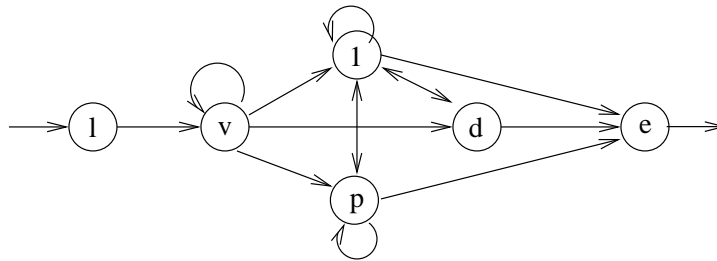


# Tik-61.231 Principles of Pattern Recognition

Answers to exercise 7: 12.11.2001

- A grammar  $G = (V_T, V_N, P, S)$  consists of the following entities:  $V_T$ , a set of terminal symbols  
 $V_N$ , a set of nonterminal symbols  
 $P$ , a set of production rules  
 $S$ , a set of starting symbols

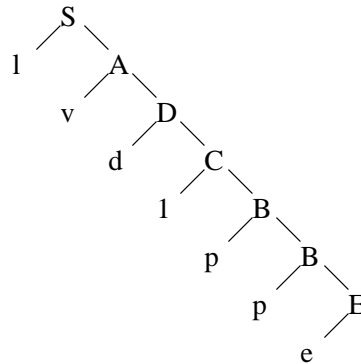
Lets use the notations  $l$  (locomotive),  $v$  (luggage van),  $p$  (passenger coach),  $1$  (first class coach),  $d$  (dining car, interpreted as a special case of a passenger coach) and  $e$  (the end of the train, an empty symbol). Now we can draw a syntax diagram for the acceptable train structures: The BNF (Backus-



Naur-Form or Backus normal form) uses the symbols  $:=$  as  $\rightarrow$ ,  $|$  for 'or' and  $\langle \dots \rangle$  for nonterminal symbols. The following grammar describes all possible trains:

$$\begin{aligned}
 V_T &= \{l, v, p, 1, d, e\} \\
 V_N &= \{S, A, B, C, D, E\} \\
 P &= \left\{ \begin{array}{l}
 \langle S \rangle := l \langle A \rangle \\
 \langle A \rangle := v(\langle A \rangle \mid \langle B \rangle \mid \langle C \rangle \mid \langle D \rangle) \\
 \langle B \rangle := p(\langle B \rangle \mid \langle C \rangle \mid \langle E \rangle) \\
 \langle C \rangle := 1(\langle B \rangle \mid \langle C \rangle \mid \langle D \rangle \mid \langle E \rangle) \\
 \langle D \rangle := d(\langle C \rangle \mid \langle D \rangle \mid \langle E \rangle) \\
 \langle E \rangle := e
 \end{array} \right\}
 \end{aligned}$$

Now the parse tree for the example train  $\{l, v, d, 1, p, p\}$  is

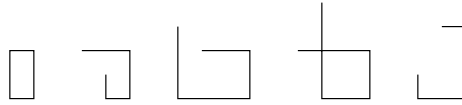


2. Now to check if the grammar accepts a train of the form  $\{l, v, v, d, p, 1\}$ . From the syntax diagram in the previous question we can instantly see, that the symbol  $d$  cannot be followed by  $p$ . Thus the grammar does not accept the train.
3. The grammar is

$$\begin{aligned}
V_T &= \{o, a, \neg, +\} \\
V_N &= \{Square, Side1, Side2, Side3, Side4\} \\
P &= \left\{ \begin{array}{l} Square \rightarrow Side1 + Side2 + Side3 + Side4 \\ Side1 \rightarrow o | Side1 + o \\ Side2 \rightarrow a | Side2 + a \\ Side3 \rightarrow \neg o | Side3 + \neg o \\ Side4 \rightarrow \neg a | Side4 + \neg a \end{array} \right\} \\
S &= \{Square\} \\
\Rightarrow L &= \{(o|o^{n_1}) + (a|a^{n_2}) + (\neg o|(\neg o)^{n_3}) + (\neg a|(\neg a)^{n_4})\}
\end{aligned}$$

Where  $n_1, n_2, n_3, n_4 \geq 1$ ,  $|$  stands for 'or' and  $a^n = a + a + a + \dots + a$  ( $n$   $a$ 's).

- a) Yes, the grammar produces squares if  $n_1 = n_2 = n_3 = n_4$ .
- b) The grammar can produce several other structures, since the only constrained factor is the order of the turns; for example (when  $n_1 = n_2 = n_3 = n_4$  doesn't hold)



- c) The grammar can be made to produce only squares by constraining the length of each vertice to be equal,

$$L = \{o^n + a^n + (\neg o)^n + (\neg a)^n | n \geq 1\}$$

Thus the grammar is

$$\begin{aligned}
V_T &= \{o, a, \neg, +\} \\
V_N &= \{N, A, V, Y\} \\
P &= \left\{ \begin{array}{l} N \rightarrow o + N + A + V + Y | o + A + V + Y \\ Y + A \rightarrow A + Y \text{ (correct order)} \\ V + A \rightarrow A + V \text{ (correct order)} \\ Y + V \rightarrow V + Y \text{ (correct order)} \\ o + A \rightarrow o + a \text{ (upper right corner)} \\ a + V \rightarrow a + \neg o \text{ (lower right corner)} \\ \neg o + Y \rightarrow \neg o + \neg a \text{ (lower left corner)} \\ a + A \rightarrow a + a \\ \neg o + V \rightarrow \neg o + \neg o \\ \neg a + Y \rightarrow \neg a + \neg a \end{array} \right\} \\
S &= \{N\}
\end{aligned}$$

4. a) The regular EKG-form, as in *Syntactic PR, An Introduction*, R.C. Gonzales and M.G. Thomason, Addison Wesley 1998, is

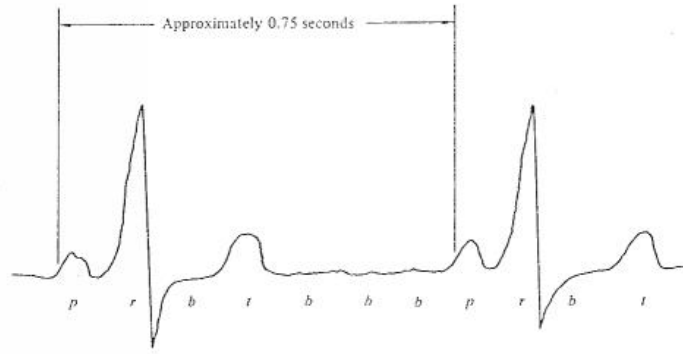
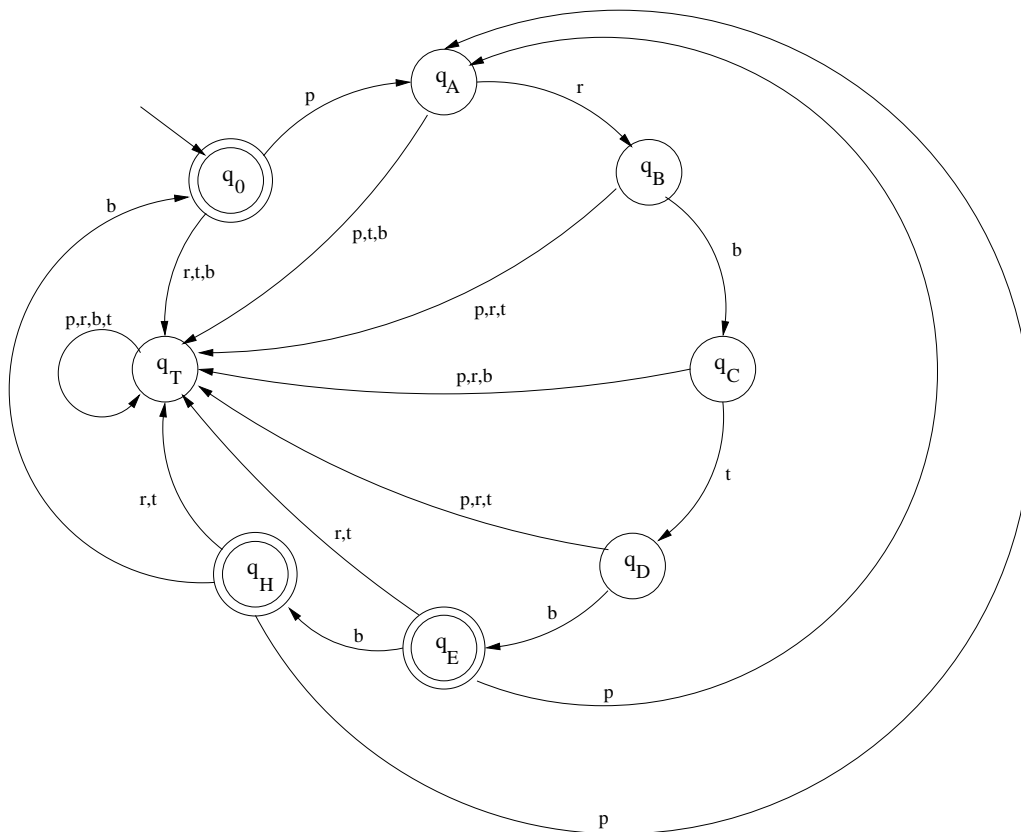


Figure 4.12. Normal human ECG.

The deterministic automaton thus becomes



- b) With the symbol sequence  $prbtbprbtbbb\dots$  the resulting state sequence is

$$q_0 \rightarrow q_A \rightarrow q_B \rightarrow q_C \rightarrow q_D \rightarrow q_E \rightarrow q_A \rightarrow q_B \rightarrow q_C \rightarrow q_D \rightarrow q_E \rightarrow q_H \rightarrow q_0$$

Now when the sequence starts again, the same route is traversed. As alarms are a result of the trap state  $q_T$ , which is not entered, no alarm is produced with this sequence.