Ensembles and probabilities Probability, Entropy, and Inference • Ensemble X is a triple $(x, \mathcal{A}_X, \mathcal{P}_X)$, where Based on David J.C. MacKay: -x is the *outcome* of random variable Information Theory, Inference and Learning Algorithms, 2003 $-\mathcal{A}_X = \{a_1, a_2, \dots, a_I\}$ are the possible values for x Chapter 2 - $\mathcal{P}_X = \{p_1, p_2, \dots, p_I\}$ are the *probabilities* of outcomes $P(x = a_i) = p_i$ $-p_i \ge 0$ $-\sum_{a_i \in \mathcal{A}_X}^{n} P(x=a_i) = 1$ Juha Raitio juha.raitio@iki.fi • $P(x = a_i)$ may be written as $P(a_i)$ or P(x)5th February 2004 • Probability of a subset T of \mathcal{A}_r $P(T) = P(x \in T) = \sum_{a_i \in T} P(x = a_i)$ (1)HUT T-61.182 Information Theory and Machine Learning Juha Raitio 5th February 2004 Outline Joint ensembles and marginal probabilities 1. On notation of probabilities • Joint ensemble XY - Outcome is an ordered pair x, y (or xy) 2. Meaning of probability - Possible values $\mathcal{A}_X = \{a_1, a_2, \dots, a_I\}$ and $\mathcal{A}_Y = \{b_1, b_2, \dots, b_J\}$ - Joint probability P(x, y)3. Forward and inverse probabilities • Marginal probabilities 4. Probabilistic inference $P(x = a_i) \equiv \sum_{y \in \mathcal{A}_Y} P(x = a_i, y)$ (2)5. Shannon information and entropy $P(y) \equiv \sum_{x \in \mathcal{A}_X} P(y, x)$ (3)6. On convexity of functions 7. Exercises

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Inference with inverse probabilities

• Inference on parameters θ given data D and hypothesis \mathcal{H} by Bayes' theorem

$$P(\theta|D,\mathcal{H}) = \frac{P(D|\theta,\mathcal{H})P(\theta|\mathcal{H})}{P(D|\mathcal{H})},$$
(10)

where

 $P(\theta|\mathcal{H})$ is the *prior probability* for parameters $P(D|\theta, \mathcal{H})$ is the *likelihood* of the parameters given the data $P(D|\mathcal{H})$ is the *evidence* $P(\theta|D, \mathcal{H})$ is the *posterior probability* for parameters

• in written

$$posterior = \frac{likelihood \times prior}{evidence}$$

Shannon information and entropy

 $h(x = a_i) = \log_2 \frac{1}{P(x = a_i)}$

 $H(X) \equiv \sum_{x \in \mathcal{A}_X} P(x) \log \frac{1}{P(x)}$

• Shannon information content of an outcome $x = a_i$ (bits)

Decomposability of entropy

• Entropy of probability distribution $\mathbf{p} = \{p_1, p_2, \dots, p_I\}$

$$H(\mathbf{p}) = H(p_1, 1 - p_1) + (1 - p_1)H\left(\frac{p_2}{1 - p_1}, \frac{p_3}{1 - p_1}, \dots, \frac{p_I}{1 - p_1}\right)$$
(15)

• More generally

$$H(\mathbf{p}) = H[(p_1 + p_2 + \dots + p_m), (p_{m+1} + p_{m+2} + \dots + p_I)] + (p_1 + \dots + p_m)H\left(\frac{p_1}{(p_1 + \dots + p_m)}, \dots, \frac{p_m}{(p_1 + \dots + p_m)}\right) \quad (16) + (p_{m+1} + \dots + p_I)H\left(\frac{p_{m+1}}{(p_{m+1} + \dots + p_I)}, \dots, \frac{p_I}{(p_{m+1} + \dots + p_I)}\right)$$

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Relative entropy

• *Kullback-Leibler divergence* between P(x) and Q(x) over alphabet A_X

$$D_{KL}(P||Q) = \sum_{x} P(x) \log \frac{P(x)}{Q(x)}$$
(17)

• Properties of relative entropy

- Gibbs' inequality:
$$D_{KL}(P||Q) \ge 0$$
 and $D_{KL}(P||Q) = 0$, if $P = Q$
- in general $D_{KL}(P||Q) \ne D_{KL}(Q||P)$

• Joint entropy of X, Y

• *Entropy of an ensemble X* (bits)

$$H(X,Y) \equiv \sum_{xy \in \mathcal{A}_X \mathcal{A}_Y} P(x,y) \log \frac{1}{P(x,y)}$$
(14)

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(11)

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(12)

(13)

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